Section 8.8 Recursive Definitions

Arithmetic Progressions

• Recall from section 2.4 that an arithmetic progression is a sequence of the form:

$$a, a + d, a + 2d, \cdots a + nd, \cdots$$

• For example, the series:

is an arithmetic progression where a = -1 and d = 4

Arithmetic Progressions

• For the arithmetic progression:

-1, 3, 7, 11, …

- Given any element of the sequence, it is a simple matter to produce the next element in the sequence by adding 4 to it
- This rule can be used to calculate any element of the sequence except for the first element: -1

• We can then create a function to compute the *n*th element of

-1, 3, 7, 11, …

by giving two rules:

1.
$$f(n) = -1$$
 if $n = 0$

2.
$$f(n) = f(n-1) + 4$$
 if $n > 0$

Note that the definition of f uses f. This is called <u>recursion</u>, and f is a <u>recursive function</u>

• Instead of using n > 0 as a case for an argument of f, it is common to express the argument as n + 1, because if n is a natural number, then n + 1 > 0

$$f(0) = -1$$
Base case $f(n+1) = 4 + f(n)$ Recursive case

• Examples of computing using the recursive function \boldsymbol{f}

$$f(0) = -1$$

• Examples of computing using the recursive function \boldsymbol{f}

$$f(0) = -1$$

 $f(1)$

$$f(0) = -1$$

$$f(1) = 4 + f(0)$$

$$f(0) = -1$$

$$f(1) = 4 + f(0)$$

= 4 + -1

$$f(0) = -1$$

$$f(1) = 4 + f(0)$$

= 4 + -1
= 3

 \bullet Examples of computing using the recursive function f

f(2)

$$f(2) = 4 + f(1)$$

$$f(2) = 4 + f(1)$$

= 4 + 4 + f(0)

$$f(2) = 4 + f(1)$$

= 4 + 4 + f(0)
= 4 + 4 + -1

$$f(2) = 4 + f(1)$$

= 4 + 4 + f(0)
= 4 + 4 + -1
= 7

$$f(5) = 4 + f(4)$$

$$f(5) = 4 + f(4)$$

= 4 + 4 + f(3)

$$f(5) = 4 + f(4)$$

= 4 + 4 + f(3)
= 4 + 4 + 4 + f(2)

$$f(5) = 4 + f(4)$$

= 4 + 4 + f(3)
= 4 + 4 + 4 + f(2)
= 4 + 4 + 4 + 4 + f(1)

$$f(5) = 4 + f(4)$$

= 4 + 4 + f(3)
= 4 + 4 + 4 + 4 + f(2)
= 4 + 4 + 4 + 4 + 4 + f(1)
= 4 + 4 + 4 + 4 + 4 + f(0)

• Examples of computing the recursive function f

f

$$(5) = 4 + f(4)$$

= 4 + 4 + f(3)
= 4 + 4 + 4 + 4 + f(2)
= 4 + 4 + 4 + 4 + 4 + f(1)
= 4 + 4 + 4 + 4 + 4 + f(0)
= 4 + 4 + 4 + 4 + 4 + -1

$$f(5) = 4 + f(4)$$

= 4 + 4 + f(3)
= 4 + 4 + 4 + 4 + f(2)
= 4 + 4 + 4 + 4 + 4 + f(1)
= 4 + 4 + 4 + 4 + 4 + 4 + f(0)
= 4 + 4 + 4 + 4 + 4 + -1
= 19

- Example 2: Define a recursive function that computes a^n where a is a real number and n is a natural number
- The underlying sequence is

$$a^{0}, a^{1}, a^{2}, \cdots$$

- Given any number in the sequence, multiply it by *a* to get the next number in the sequence
- The first number in the sequence is 1

$$f(0) = 1$$
$$f(n+1) = a \cdot f(n)$$

$$f(4) = a \cdot f(3)$$

$$f(4) = a \cdot f(3)$$
$$= a \cdot a \cdot f(2)$$

$$f(4) = a \cdot f(3)$$
$$= a \cdot a \cdot f(2)$$
$$= a \cdot a \cdot a \cdot f(1)$$

$$f(4) = a \cdot f(3)$$

= $a \cdot a \cdot f(2)$
= $a \cdot a \cdot a \cdot a \cdot f(1)$
= $a \cdot a \cdot a \cdot a \cdot a \cdot f(0)$

$$F(4) = a \cdot f(3)$$

= $a \cdot a \cdot a \cdot f(2)$
= $a \cdot a \cdot a \cdot a \cdot f(1)$
= $a \cdot a \cdot a \cdot a \cdot a \cdot f(0)$
= $a \cdot a \cdot a \cdot a \cdot a \cdot 1$

$$f(4) = a \cdot f(3)$$

= $a \cdot a \cdot a \cdot f(2)$
= $a \cdot a \cdot a \cdot a \cdot f(1)$
= $a \cdot a \cdot a \cdot a \cdot a \cdot f(0)$
= $a \cdot a \cdot a \cdot a \cdot 1$
= a^4

- Example 3: Define a recursive function that computes $\sum_{k=0}^{n} k$ where n is a natural number
- The underlying sequence of sums is

 $0, 0 + 1, 0 + 1 + 2, \cdots$

- Given the nth sum in the sequence, add n + 1 to get the next sum in the sequence
- The first, 0th, sum in the sequence is 0

$$f(0) = 0$$

 $f(n+1) = n + 1 + f(n)$

$$f(4) = 4 + f(3)$$

$$f(4) = 4 + f(3)$$

= 4 + 3 + f(2)

$$f(4) = 4 + f(3)$$

= 4 + 3 + f(2)
= 4 + 3 + 2 + f(1)

$$f(4) = 4 + f(3)$$

= 4 + 3 + f(2)
= 4 + 3 + 2 + f(1)
= 4 + 3 + 2 + 1 + f(0)
• Example 3: Define a recursive function that computes $\sum_{k=0}^{n} k$ where n is a natural number

$$f(4) = 4 + f(3)$$

= 4 + 3 + f(2)
= 4 + 3 + 2 + f(1)
= 4 + 3 + 2 + 1 + f(0)
= 4 + 3 + 2 + 1 + 0

• Example 3: Define a recursive function that computes $\sum_{k=0}^{n} k$ where n is a natural number

$$f(4) = 4 + f(3)$$

= 4 + 3 + f(2)
= 4 + 3 + 2 + f(1)
= 4 + 3 + 2 + 1 + f(0)
= 4 + 3 + 2 + 1 + 0
= 10

- Example 4: Define a recursive function that computes n factorial: $n! = 1 \cdot 2 \cdot ... \cdot n$ where n is a natural number
- Note that 0! = 1
- The underlying sequence of products is:

1, $1 \cdot 1$, $1 \cdot 1 \cdot 2$, $1 \cdot 1 \cdot 2 \cdot 3$, ...

- Given the nth product in the sequence, multiply it by n + 1 to get the next product in the sequence
- The first, 0th, product in the sequence is 1

$$f(0) = 1$$
$$f(n+1) = f(n) \cdot (n+1)$$

$$f(4) = f(3) \cdot 4$$

$$f(4) = f(3) \cdot 4$$
$$= f(2) \cdot 3 \cdot 4$$

$$f(4) = f(3) \cdot 4$$

= $f(2) \cdot 3 \cdot 4$
= $f(1) \cdot 2 \cdot 3 \cdot 4$

$$f(4) = f(3) \cdot 4$$

= $f(2) \cdot 3 \cdot 4$
= $f(1) \cdot 2 \cdot 3 \cdot 4$
= $f(0) \cdot 1 \cdot 2 \cdot 3 \cdot 4$

$$f(4) = f(3) \cdot 4$$

= f(2) \cdot 3 \cdot 4
= f(1) \cdot 2 \cdot 3 \cdot 4
= f(0) \cdot 1 \cdot 2 \cdot 3 \cdot 4
= 1 \cdot 1 \cdot 2 \cdot 3 \cdot 4

$$f(4) = f(3) \cdot 4$$

= $f(2) \cdot 3 \cdot 4$
= $f(1) \cdot 2 \cdot 3 \cdot 4$
= $f(0) \cdot 1 \cdot 2 \cdot 3 \cdot 4$
= $1 \cdot 1 \cdot 2 \cdot 3 \cdot 4$
= 24

• The Fibonacci sequence is

0, 1, 1, 2, 3, 5, 8, 13, …

• The Fibonacci sequence is

0, 1, 1, 2, 3, 5, 8, 13, …

The first two numbers of the sequence are 0 and 1. Each other number in the sequence is the sum of its two previous numbers in the sequence

$$f(0) = 0$$

$$f(1) = 1$$

$$f(n+2) = f(n) + f(n+1)$$

• The Fibonacci function computes values in the Fibonacci sequence

f(4) = f(2) + f(3)

• The Fibonacci function computes values in the Fibonacci sequence

f(4) = f(2) + f(3)= f(0) + f(1) + f(3)

$$f(4) = f(2) + f(3)$$

= f(0) + f(1) + f(3)
= 0 + f(1) + f(3)

$$f(4) = f(2) + f(3)$$

= f(0) + f(1) + f(3)
= 0 + f(1) + f(3)
= 0 + 1 + f(3)

$$f(4) = f(2) + f(3)$$

= f(0) + f(1) + f(3)
= 0 + f(1) + f(3)
= 0 + 1 + f(3)
= 0 + 1 + f(1) + f(2)

$$f(4) = f(2) + f(3)$$

= f(0) + f(1) + f(3)
= 0 + f(1) + f(3)
= 0 + 1 + f(3)
= 0 + 1 + f(1) + f(2)
= 0 + 1 + 1 + f(2)

$$f(4) = f(2) + f(3)$$

= f(0) + f(1) + f(3)
= 0 + f(1) + f(3)
= 0 + 1 + f(3)
= 0 + 1 + f(1) + f(2)
= 0 + 1 + 1 + f(2)
= 0 + 1 + 1 + f(0) + f(1)

The Fibonacci function computes values in the Fibonacci sequence

f(4) = f(2) + f(3)= f(0) + f(1) + f(3) = 0 + f(1) + f(3) = 0 + 1 + f(3) = 0 + 1 + f(1) + f(2) = 0 + 1 + 1 + f(2) = 0 + 1 + 1 + f(0) + f(1) = 0 + 1 + 1 + 0 + f(1)

The Fibonacci function computes values in the Fibonacci sequence

f(4) = f(2) + f(3)= f(0) + f(1) + f(3)= 0 + f(1) + f(3)= 0 + 1 + f(3)= 0 + 1 + f(1) + f(2)= 0 + 1 + 1 + f(2)= 0 + 1 + 1 + f(0) + f(1)= 0 + 1 + 1 + 0 + f(1)= 0 + 1 + 1 + 0 + 1

The Fibonacci function computes values in the Fibonacci sequence

f(4) = f(2) + f(3)= f(0) + f(1) + f(3)= 0 + f(1) + f(3)= 0 + 1 + f(3)= 0 + 1 + f(1) + f(2)= 0 + 1 + 1 + f(2)= 0 + 1 + 1 + f(0) + f(1)= 0 + 1 + 1 + 0 + f(1)= 0 + 1 + 1 + 0 + 1= 3

- Recursively defined sets use recursion to specify the elements in a set
 - 1. Base elements of the set are explicitly defined
 - 2. A recursive rule is given to define additional elements in the set

• Recursively defined sets are also known as inductively defined sets

- Example: A recursive definition of *N*, the set of natural numbers
 - 1. $0 \in N$
 - 2. If $x \in N$, then $x + 1 \in N$
 - 3. Nothing else is in **N**

• Example 5: A recursive definition of a subset *S* of natural numbers

- 1. $3 \in S$
- 2. If $x \in S$ and $y \in S$, then $x + y \in S$
- 3. Nothing else is in S

• Example: A recursive definition of the set of properly nested parentheses, *P*

- 1. () $\in P$
- 2. If $u \in P$ and $v \in P$, then $(u) \in P$ and $uv \in P$
- 3. Nothing else is in *P*

Recursively Defined Set of Binary Strings

• The set of binary strings of any finite, non-negative length, B^* has a recursive definition

- 1. $\lambda \in B^*$ where λ is the empty string
- 2. If $s \in B^*$, then $s0 \in B^*$ and $s1 \in B^*$
- 3. Nothing else is in B^*

Recursively Defined Set of Strings

- Some members of B^*
 - $\lambda \in B^*$
 - $\lambda 0 = 0$, so $0 \in B^*$
 - $\lambda 1 = 1$, so $1 \in B^*$
 - $00 \in B^*$
 - $10 \in B^*$
 - $01 \in B^*$
 - $11 \in B^*$

Recursive Functions on Binary Strings

- Let $|\cdot|: B^* \to N$ be a function that recursively computes the length of a string
 - Note that $|\cdot|$ is a function that is called by replacing the dot with an argument

$$|\lambda| = 0$$

 $|s0| = 1 + |s|$
 $|s1| = 1 + |s|$

|0110| = 1 + |011|

|0110| = 1 + |011|= 1 + 1 + |01|

|0110| = 1 + |011|= 1 + 1 + |01| = 1 + 1 + 1 + |1|

|0110| = 1 + |011|= 1 + 1 + |01| = 1 + 1 + 1 + |0| = 1 + 1 + 1 + 1 + |\lambda|

|0110| = 1 + |011|= 1 + 1 + |01| = 1 + 1 + 1 + |1| = 1 + 1 + 1 + 1 + | λ | = 1 + 1 + 1 + 1 + 0

|0110| = 1 + |011|= 1 + 1 + |01| = 1 + 1 + 1 + |1| = 1 + 1 + 1 + 1 + |\lambda| = 1 + 1 + 1 + 1 + 0 = 4
Recursive Functions on Strings

• Compare the recursive definition of B^* to the recursive definition of $|\cdot|: B^* \to N$

$\lambda \in B^*$	$ \lambda = 0$
$s0 \in B^*$ if $s \in B^*$	s0 = 1 + s
$s1 \in B^*$ if $s \in B^*$	s1 = 1 + s

• Example: Define a recursive function that counts the number of digits in a natural number

• First attempt

length(0) = 1length(n + 1) = ???

- Consider a different definition of the natural numbers
 - 1. If $d \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, then $d \in N$
 - 2. If $n \in N$ and $d \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, then $10n + d \in N$

- Consider a different definition of the natural numbers
 - 1. If $d \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, then $d \in N$
 - 2. If $n \in N$ and $d \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, then $10n + d \in N$

Examples:

 $4 \in N$ $45 \in N$ $451 \in N$

- Consider a different definition of the natural numbers
 - 1. If $d \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, then $d \in N$
 - 2. If $n \in N$ and $d \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, then $10n + d \in N$

length(0) = 1 length(1) = 1 length(2) = 1 length(3) = 1 length(4) = 1length(5) = 1 length(6) = 1 length(7) = 1 length(8) = 1 length(9) = 1length(10n + d) = 1 + length(n)

 $length(0) = 1 \quad length(1) = 1 \quad length(2) = 1 \quad length(3) = 1 \quad length(4) = 1$ $length(5) = 1 \quad length(6) = 1 \quad length(7) = 1 \quad length(8) = 1 \quad length(9) = 1$ length(10n + d) = 1 + length(n)

length(451)length(451)

length(0) = 1 length(1) = 1 length(2) = 1 length(3) = 1 length(4) = 1length(5) = 1 length(6) = 1 length(7) = 1 length(8) = 1 length(9) = 1length(10n + d) = 1 + length(n)

length(451) = 1 + length(45)length(451)

length(0) = 1 length(1) = 1 length(2) = 1 length(3) = 1 length(4) = 1length(5) = 1 length(6) = 1 length(7) = 1 length(8) = 1 length(9) = 1length(10n + d) = 1 + length(n)

length(451) = 1 + length(45)length(451) = 1 + 1 + length(4)

length(0) = 1 length(1) = 1 length(2) = 1 length(3) = 1 length(4) = 1length(5) = 1 length(6) = 1 length(7) = 1 length(8) = 1 length(9) = 1length(10n + d) = 1 + length(n)

length(451) = 1 + length(45)length(451) = 1 + 1 + length(4)= 1 + 1 + 1

length(0) = 1 length(1) = 1 length(2) = 1 length(3) = 1 length(4) = 1length(5) = 1 length(6) = 1 length(7) = 1 length(8) = 1 length(9) = 1length(10n + d) = 1 + length(n)

```
length(451) = 1 + length(45)
length(451) = 1 + 1 + length(4)
= 1 + 1 + 1
= 3
```

length(0) = 1 length(1) = 1 length(2) = 1 length(3) = 1 length(4) = 1length(5) = 1 length(6) = 1 length(7) = 1 length(8) = 1 length(9) = 1length(10n + d) = 1 + length(n)

Note that $n = \left\lfloor \frac{(10n+d)}{10} \right\rfloor$

- Rewrite as a function in pseudocode
 - Name: length
 - Input: a natural number *n*
 - Output: The number of digits in *n*

```
if n <= 9
  return 1
else
  return 1 + length(floor(n/10))
end-if</pre>
```