

# Compound Propositions - Conditional

- Also known as "implication"
- $p \rightarrow q$
- Truth table:

$p$	$q$	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

- "if  $p$  then  $q$ ", " $p$  implies  $q$ ", " $p$  only if  $q$ ", " $p$  is sufficient for  $q$ ", " $q$  is necessary for  $p$ ", " $q$  unless  $\neg p$ "

# Compound Propositions - Conditional

- Assume  $p$  represents “You will get 100% on the final exam”
- and  $q$  represents “You will get an 'A' for the course”
- Then  $p \rightarrow q$  represents
  - “If you will get 100% on the final exam then you will get an 'A' for the course.”
  - Note: nothing is guaranteed if you do not get 100% on the final; you could still get an 'A'.

# Compound Propositions - Conditional

- Assume  $p$  represents “John wakes up early”
- and  $q$  represents “John will be late for the meeting”
- Then  $(\neg p) \rightarrow q$  represents
  - "If John does not wake up early, then John will be late for the meeting"

# Converse, Contrapositive, and Inverse

- The converse of  $p \rightarrow q$  is  $q \rightarrow p$
- The contrapositive of  $p \rightarrow q$  is  $\neg q \rightarrow \neg p$
- The inverse of  $p \rightarrow q$  is  $\neg p \rightarrow \neg q$

# Compound Propositions - Biconditional

- Also known as "bi-implication" or "double implication"
- $p \leftrightarrow q$
- Truth table:

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

- " $p$  if and only if  $q$ "

# Truth Tables for Compound Propositions

- Example:  $(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
- Since there are two variables,  $p$  and  $q$ , create a table with four rows and columns for sub-propositions

$p$	$q$	$\neg q$	$p \leftrightarrow q$	$p \vee \neg q$	$(p \leftrightarrow q) \rightarrow (p \vee \neg q)$

# Truth Tables for Compound Propositions

- Example:  $(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
- List all possible combinations of truth values for  $p$  and  $q$

$p$	$q$	$\neg q$	$p \leftrightarrow q$	$p \vee \neg q$	$(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
T	T				
T	F				
F	T				
F	F				

# Truth Tables for Compound Propositions

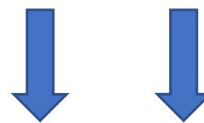
- Example:  $(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
- Complete the  $\neg q$  column using the  $q$  column



$p$	$q$	$\neg q$	$p \leftrightarrow q$	$p \vee \neg q$	$(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
T	T	F			
T	F	T			
F	T	F			
F	F	T			

# Truth Tables for Compound Propositions

- Example:  $(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
- Complete the  $p \leftrightarrow q$  column using the  $p$  and  $q$  columns



$p$	$q$	$\neg q$	$p \leftrightarrow q$	$p \vee \neg q$	$(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
T	T	F	T		
T	F	T	F		
F	T	F	F		
F	F	T	T		

# Truth Tables for Compound Propositions

- Example:  $(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
- Complete the  $p \vee \neg q$  column using the  $p$  and  $\neg q$  columns



$p$	$q$	$\neg q$	$p \leftrightarrow q$	$p \vee \neg q$	$(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
T	T	F	T	T	
T	F	T	F	T	
F	T	F	F	F	
F	F	T	T	T	

# Truth Tables for Compound Propositions

- Example:  $(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
- Complete the  $(p \leftrightarrow q) \rightarrow (p \vee \neg q)$  column using the  $p \leftrightarrow q$  and  $p \vee \neg q$  columns



$p$	$q$	$\neg q$	$p \leftrightarrow q$	$p \vee \neg q$	$(p \leftrightarrow q) \rightarrow (p \vee \neg q)$
T	T	F	T	T	T
T	F	T	F	T	T
F	T	F	F	F	T
F	F	T	T	T	T

# Example

- Create a truth table for  $(p \wedge q) \vee (r \wedge s)$

# Example

# Example

# Example

$p$	$q$	$r$	$s$	$(p \wedge q)$	$(r \wedge s)$	$(p \wedge q) \vee (r \wedge s)$
T	T	T	T			
T	T	T	F			
T	T	F	T			
T	T	F	F			
T	F	T	T			
T	F	T	F			
T	F	F	T			
T	F	F	F			
F	T	T	T			
F	T	T	F			
F	T	F	T			
F	T	F	F			
F	F	T	T			
F	F	T	F			
F	F	F	T			
F	F	F	F			

# Example

$p$	$q$	$r$	$s$	$(p \wedge q)$	$(r \wedge s)$	$(p \wedge q) \vee (r \wedge s)$
T	T	T	T	T		
T	T	T	F	T		
T	T	F	T	T		
T	T	F	F	T		
T	F	T	T	F		
T	F	T	F	F		
T	F	F	T	F		
T	F	F	F	F		
F	T	T	T	F		
F	T	T	F	F		
F	T	F	T	F		
F	T	F	F	F		
F	F	T	T	F		
F	F	T	F	F		
F	F	F	T	F		
F	F	F	F	F		

# Example

$p$	$q$	$r$	$s$	$(p \wedge q)$	$(r \wedge s)$	$(p \wedge q) \vee (r \wedge s)$
T	T	T	T	T	T	
T	T	T	F	T	F	
T	T	F	T	T	F	
T	T	F	F	T	F	
T	F	T	T	F	T	
T	F	T	F	F	F	
T	F	F	T	F	F	
T	F	F	F	F	F	
F	T	T	T	F	T	
F	T	T	F	F	F	
F	T	F	T	F	F	
F	T	F	F	F	F	
F	F	T	T	F	T	
F	F	T	F	F	F	
F	F	F	T	F	F	
F	F	F	F	F	F	

# Example

$p$	$q$	$r$	$s$	$(p \wedge q)$	$(r \wedge s)$	$(p \wedge q) \vee (r \wedge s)$
T	T	T	T	T	T	T
T	T	T	F	T	F	T
T	T	F	T	T	F	T
T	T	F	F	T	F	T
T	F	T	T	F	T	T
T	F	T	F	F	F	F
T	F	F	T	F	F	F
T	F	F	F	F	F	F
F	T	T	T	F	T	T
F	T	T	F	F	F	F
F	T	F	T	F	F	F
F	T	F	F	F	F	F
F	F	T	T	F	T	T
F	F	T	F	F	F	F
F	F	F	T	F	F	F
F	F	F	F	F	F	F

# Tautologies, Contradictions, and Contingencies

- If a proposition is true for all values of its variables, then it is a tautology.
  - $p \vee \neg p$

$p$	$\neg p$	$p \vee \neg p$
T	F	T
F	T	T

# Tautologies, Contradictions, and Contingencies

- If a proposition is false for all values of its variables, then it is a contradiction
  - $p \wedge \neg p$

$p$	$\neg p$	$p \wedge \neg p$
T	F	F
F	T	F

# Tautologies, Contradictions, and Contingencies

- If a proposition is neither a tautology nor a contradiction, then it is called a contingency
  - $p \wedge q$

$p$	$q$	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

# Logical Equivalences

- Two propositions,  $p$  and  $q$ , are logically equivalent if  $p \leftrightarrow q$  is a tautology.
- The notation  $p \equiv q$  denotes that  $p$  and  $q$  are logically equivalent.
  - Sometimes  $p \Leftrightarrow q$  is used instead of  $p \equiv q$
- Note that  $p \equiv q$  is not a proposition. It is a statement about  $p \leftrightarrow q$  being a tautology

# Logical Equivalences

- Example:  $p \rightarrow q$  and  $\neg p \vee q$  are logically equivalent

$p$	$q$	$p \rightarrow q$	$\neg p$	$\neg p \vee q$	$p \rightarrow q \leftrightarrow \neg p \vee q$
T	T	T	F	T	T
T	F	F	F	F	T
F	T	T	T	T	T
F	F	T	T	T	T

# De Morgan's Laws

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

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$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

$p$	$q$	$p \wedge q$	$\neg(p \wedge q)$	$\neg p$	$\neg q$	$\neg p \vee \neg q$	$\neg(p \wedge q) \leftrightarrow \neg p \vee \neg q$
T	T	T	F	F	F	F	T
T	F	F	T	F	T	T	T
F	T	F	T	T	F	T	T
F	F	F	T	T	T	T	T

# De Morgan's Laws

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

$p$	$q$	$p \vee q$	$\neg(p \vee q)$	$\neg p$	$\neg q$	$\neg p \wedge \neg q$	$\neg(p \vee q) \leftrightarrow \neg p \wedge \neg q$
T	T	T	F	F	F	F	T
T	F	T	F	F	T	F	T
F	T	T	F	T	F	F	T
F	F	F	T	T	T	T	T

# De Morgan's Laws

- De Morgan's laws can be applied to cases with compound propositions as well as propositional variables. Substitute the same compound proposition for all occurrences of the same variable

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

Substitute  $s \rightarrow t$  for  $p$  and  $p \wedge q$  for  $q$

$$\neg((s \rightarrow t) \vee (p \wedge q)) \equiv \neg(s \rightarrow t) \wedge \neg(p \wedge q)$$

# De Morgan's Laws and Natural Language

- Use De Morgan's laws to express the negation of:  
"Miguel has a cellphone and he has a laptop computer"
- Let  $p$  represent "Miguel has a cellphone"
- Let  $q$  represent "Miguel has a laptop computer"
- $\neg(p \wedge q) \equiv \neg p \vee \neg q$
- "Miguel does not have a cellphone or he does not have a laptop computer"

# Examples of Using De Morgan's Laws

- $\neg((p \vee (r \rightarrow q)) \wedge (r \vee s))$

# Examples of Using De Morgan's Laws

- $\neg((p \vee (r \rightarrow q)) \wedge (r \vee s))$
- $\neg(p \vee (r \rightarrow q)) \vee \neg(r \vee s)$

# Examples of Using De Morgan's Laws

- $\neg((p \vee s) \vee r) \wedge \neg(p \vee q)$

# Examples of Using De Morgan's Laws

- $\neg((p \vee s) \vee r) \wedge \neg(p \vee q)$
- $\neg(((p \vee s) \vee r) \vee (p \vee q))$

# Examples of Using De Morgan's Laws

- $\neg((\neg r \rightarrow s) \wedge (\neg p \vee r)) \vee \neg((s \wedge \neg t) \vee (r \rightarrow p))$

# Examples of Using De Morgan's Laws

- $\neg((\neg r \rightarrow s) \wedge (\neg p \vee r)) \vee \neg((s \wedge \neg t) \vee (r \rightarrow p))$
- $\neg(((\neg r \rightarrow s) \wedge (\neg p \vee r)) \wedge ((s \wedge \neg t) \vee (r \rightarrow p)))$