

# Logical Equivalence in Predicate Logic

- If two statements in predicate logic are logically equivalent, then they have the same truth value regardless of the meaning of their predicates.
- Example:

$$\forall x(P(x) \wedge Q(x)) \equiv \forall xP(x) \wedge \forall xQ(x)$$

# Logical Equivalence in Predicate Logic

- Note that it is NOT the case that:

$$\forall x(P(x) \vee Q(x)) \equiv \forall xP(x) \vee \forall xQ(x)$$

- We can show that they are not logically equivalent by giving a domain of discourse and interpretations of  $P$  and  $Q$  such that the two statements have different truth values

# Logical Equivalence in Predicate Logic

- Note that it is NOT the case that:

$$\forall x(P(x) \vee Q(x)) \equiv \forall xP(x) \vee \forall xQ(x)$$

- Let the domain of discourse be the integers, let  $P(x)$  be true exactly when  $x$  is an even number, and let  $Q(x)$  be true exactly when  $x$  is an odd number.
- With that interpretation,  $\forall x(P(x) \vee Q(x))$  is true
- But  $\forall xP(x) \vee \forall xQ(x)$  is false

# Negated Quantifiers

- "x is a parent"

$$\exists y P(x, y)$$

- Note that "x is not a parent" is the same as saying, "For any person y, x is not the parent of y":

$$\neg \exists y P(x, y) \equiv \forall y \neg P(x, y)$$

- In general, the negation of an existentially quantified statement is the same as universally quantifying the negation of the inner statement

# Negated Quantifiers

- “Everybody loves  $z$ ”

$$\forall x L(x, z)$$

- The negation of “Everybody loves  $z$ ” means that there is a person that does not love  $x$ :

$$\neg \forall x L(x, z) \equiv \exists x \neg L(x, z)$$

# De Morgan's Laws for Quantified Statements

$$\neg \forall x P(x) \equiv \exists x \neg P(x)$$

$$\neg \exists x P(x) \equiv \forall x \neg P(x)$$

- As with other laws  $P(x)$  is a template that can be replaced with more complex statements. Example:

$$\neg \forall x (Q(x) \wedge R(x, y)) \equiv \exists x \neg (Q(x) \wedge R(x, y))$$

# De Morgan's Laws for Quantified Statements

- De Morgan's laws for conjunction and disjunction also can be used in predicate logic:

$$\begin{aligned}\neg \forall x (Q(x) \wedge R(x, y)) &\equiv \exists x \neg (Q(x) \wedge R(x, y)) \\ &\equiv \exists x (\neg Q(x) \vee \neg R(x, y))\end{aligned}$$